# CS61B Lecture #36

**Today:**

- A Brief Side Trip: Enumeration types.

- *DSIJ,* Chapter 10, *HFJ,* pp. 489–516.

  - Threads

  - Communication between threads

  - Synchronization

  - Mailboxes

# Side Trip into Java: Enumeration Types

- Problem: Need a type to represent something that has a few, named, discrete values.

- In the purest form, the only necessary operations are == and !=; the only property of a value of the type is that it differs from all others.

- In older versions of Java, used named integer constants:

```
interface Pieces {
   int BLACK_PIECE = 0,     // Fields in interfaces are static final.
        BLACK_KING = 1,
        WHITE_PIECE = 2,
        WHITE_KING = 3,
        EMPTY = 4;
}
```

- C and C++ provide *enumeration types* as a shorthand, with syntax like this:

```
enum Piece { BLACK_PIECE, BLACK_KING, WHITE_PIECE, WHITE_KING, EMPTY };
```

- But since all these values are basically **ints**, accidents can happen.

# Enum Types in Java

- New version of Java allows syntax like that of C or C++, but with more guarantees:

  ```java
  public enum Piece {
      BLACK_PIECE, BLACK_KING, WHITE_PIECE, WHITE_KING, EMPTY
  }
  ```

- Defines `Piece` as a new reference type, a special kind of class type.

- The names `BLACK_PIECE`, etc., are static, final *enumeration constants* (or *enumerals*) of type `PIECE`.

- They are automatically initialized, and are the only values of the enumeration type that exist (illegal to use **new** to create an enum value.)

- Can safely use `==`, and also `switch` statements:

  ```java
  boolean isKing(Piece p) {
      switch (p) {
          case BLACK_KING: case WHITE_KING: return true;
          default: return false;
      }
  }
  ```

# Making Enumerals Available Elsewhere

- Enumerals like `BLACK_PIECE` are static members of a class, not classes.

- Therefore, unlike C or C++, their declarations are not automatically visible outside the enumeration class definition.

- So, in other classes, must write `Piece.BLACK_PIECE`, which can get annoying.

- However, with version 1.5, Java has *static imports:* to import all static definitions of class `checkers.Piece` (including enumerals), you write

      import static checkers.Piece.*;

  among the import clauses.

- Alas, *cannot* use this for enum classes in the anonymous package.

# Operations on Enum Types

- Order of declaration of enumeration constants significant: `.ordinal()` gives the position (numbering from 0) of an enumeration value. Thus, `Piece.BLACK_KING.ordinal()` is 1.

- The array `Piece.values()` gives all the possible values of the type. Thus, you can write:

```
for (Piece p : Piece.values())
    System.out.printf("Piece value #%d is %s%n", p.ordinal(), p);
```

- The static function `Piece.valueOf` converts a String into a value of type `Piece`. So `Piece.valueOf("EMPTY") == EMPTY`.

# Fancy Enum Types

- Enums are classes. You can define all the extra fields, methods, and constructors you want.

- Constructors are used only in creating enumeration constants. The constructor arguments follow the constant name:

```java
enum Piece {
    BLACK_PIECE(BLACK, false, "b"), BLACK_KING(BLACK, true, "B"),
    WHITE_PIECE(WHITE, false, "w"), WHITE_KING(WHITE, true, "W"),
    EMPTY(null, false, " ");

    private final Side color;
    private final boolean isKing;
    private final String textName;

    Piece(Side color, boolean isKing, String textName) {
        this.color = color; this.isKing = isKing; this.textName = textName;
    }

    Side color() { return color; }
    boolean isKing() { return isKing; }
    String textName() { return textName; }
}
```

# Threads

- So far, all our programs consist of single sequence of instructions.

- Each such sequence is called a *thread* (for "thread of control") in Java.

- Java supports programs containing *multiple* threads, which (conceptually) run concurrently.

- Actually, on a uniprocessor, only one thread at a time actually runs, while others wait, but this is largely invisible.

- To allow program access to threads, Java provides the type `Thread` in `java.lang`. Each `Thread` contains information about, and controls, one thread.

- Simultaneous access to data from two threads can cause chaos, so are also constructs for controlled communication, allowing threads to *lock* objects, to *wait* to be notified of events, and to *interrupt* other threads.

# But Why?

- Typical Java programs always have $> 1$ thread: besides the main program, others clean up garbage objects, receive signals, update the display, other stuff.

- When programs deal with asynchronous events, is sometimes convenient to organize into subprograms, one for each independent, related sequence of events.

- Threads allow us to insulate one such subprogram from another.

- GUIs often organized like this: application is doing some computation or I/O, another thread waits for mouse clicks (like 'Stop'), another pays attention to updating the screen as needed.

- Large servers like search engines may be organized this way, with one thread per request.

- And, of course, sometimes we *do* have a real multiprocessor.

# Java Mechanics

- To specify the actions "walking" and "chewing gum":

```java
class Chewer1 implements Runnable {
   public void run()
     { while (true) ChewGum(); }
}
class Walker1 implements Runnable {
   public void run()
     { while (true) Walk(); }
}
```

```java
// Walk and chew gum
Thread chomp
  = new Thread(new
Chewer1());
Thread clomp
  = new Thread(new
Walker1());
chomp.start(); clomp.start();
```

- Concise Alternative (uses fact that Thread implements Runnable):

```java
class Chewer2 extends Thread {
  public void run()
    { while (true) ChewGum(); }
}
class Walker2 extends Thread {
   public void run()
     { while (true) Walk(); }
}
```

```java
Thread chomp = new Chewer2(),
        clomp = new Walker2();
chomp.start();
clomp.start();
```

# Avoiding Interference

- When one thread has data for another, one must wait for the other to be ready.

- Likewise, if two threads use the same data structure, generally only one should modify it at a time; other must wait.

- E.g., what would happen if two threads simultaneously inserted an item into a linked list at the same point in the list?

- A: Both could conceivably execute

  ```
  p.next = new ListCell(x, p.next);
  ```

  with the *same* values of `p` and `p.next`; one insertion is lost.

- Can arrange for only one thread at a time to execute a method on a particular object with either of the following equivalent definitions:

  ```
  void f(...) {
    synchronized (this) {
      body of f
    }
  }
  ```

  ```
  synchronized void f(...) {
    body of f
  }
  ```

# Communicating the Hard Way

- Communicating data is tricky: the faster party must wait for the slower.

- Obvious approaches for sending data from thread to thread don't work:

```
class DataExchanger {
    Object value = null;
    Object receive() {
        Object r; r = null;
        while (r == null)
           { r = value; }
        value = null;
        return r;
    }
    void deposit(Object data) {
        while (value != null) { }
        value = data;
    }
 }
```

```
DataExchanger exchanger
  = new DataExchanger();

---------------------------------

// thread1 sends to thread2 with
exchanger.deposit("Hello!");

---------------------------------

// thread2 receives from thread1 with
msg = (String) exchanger.receive();
```

- BAD: One thread can monopolize machine while waiting; two threads executing `deposit` or `receive` simultaneously cause chaos.

# Primitive Java Facilities

- `wait` method on Object makes thread wait (not using processor) until notified by `notifyAll`, unlocking the Object while it waits.

- Example, `ucb.util.mailbox` has something like this (simplified):

```java
interface Mailbox {
  void deposit(Object msg) throws InterruptedException;
  Object receive() throws InterruptedException;
}

class QueuedMailbox implements Mailbox {
  private List<Object> queue = new LinkedList<Object>();

  public synchronized void deposit(Object msg) {
    queue.add(msg);
    this.notifyAll(); // Wake any waiting receivers
  }

  public synchronized Object receive() throws InterruptedException {
    while (queue.isEmpty()) wait();
    return queue.remove(0);
  }
}
```
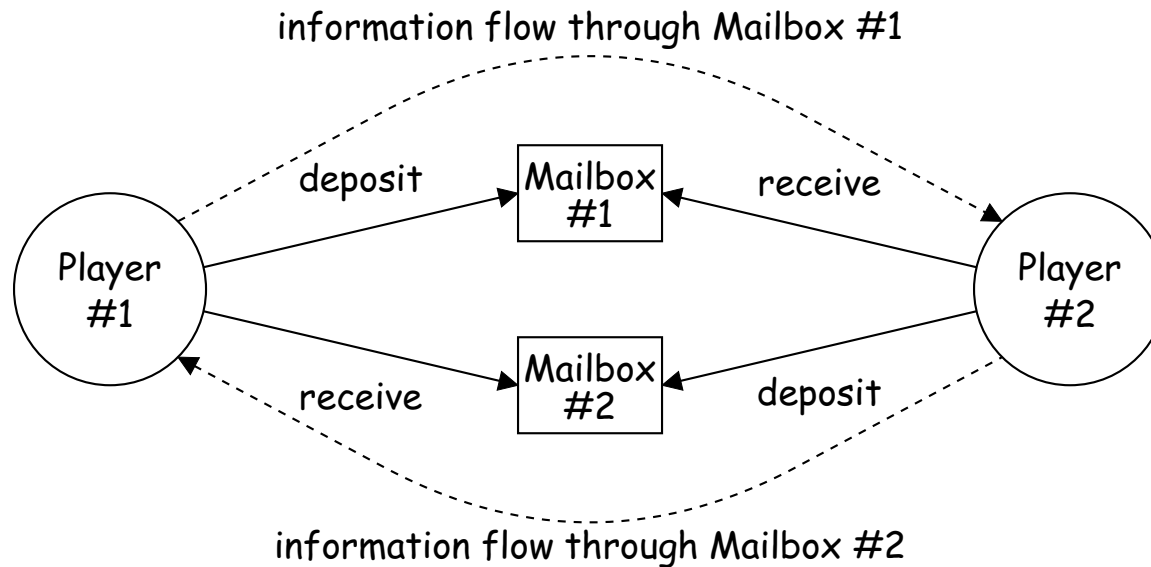
# Message-Passing Style

- Use of Java primitives very error-prone. Wait until CS162.

- Mailboxes are higher-level, and allow the following program struc-
ture:



information flow through Mailbox #1

deposit    Mailbox #1    receive

Player #1    Player #2

Mailbox #2

receive    deposit

information flow through Mailbox #2

- Where each Player is a thread that looks like this:

```
while (! gameOver()) {
  if (myMove())
    outBox.deposit(computeMyMove(lastMove));
  else
    lastMove = inBox.receive();
}
```

# More Concurrency

- Previous example can be done other ways, but mechanism is very flexible.

- E.g., suppose you want to think during opponent's move:

```
while (!gameOver()) {
  if (myMove())
    outBox.deposit(computeMyMove(lastMove));
  else {
    do {
      thinkAheadALittle();
      lastMove = inBox.receiveIfPossible();
    } while (lastMove == null);
  }
}
```

- `receiveIfPossible` (written `receive(0)` in our actual package) doesn't wait; returns `null` if no message yet, perhaps like this:

```
public synchronized Object receiveIfPossible()
     throws InterruptedException {
  if (queue.isEmpty())
    return null;
  return queue.remove(0);
}
```

# Coroutines

- A *coroutine* is a kind of synchronous thread that explicitly hands off control to other coroutines so that only one executes at a time, like Python generators. Can get similar effect with threads and mailboxes.

- Example: recursive inorder tree iterator:

```java
class TreeIterator extends Thread {
  Tree root; Mailbox r;
  TreeIterator(Tree T, Mailbox r) {
    this.root = T; this.dest = r;
  }
  public void run() {
    traverse(root);
    r.deposit(End marker);
  }
  void traverse(Tree t) {
    if (t == null) return;
    traverse(t.left);
    r.deposit(t.label);
    traverse(t.right);
  }
}

  void treeProcessor(Tree T) {
    Mailbox m = new QueuedMailbox();
    new TreeIterator(T, m).start();
    while (true) {
      Object x = m.receive();
      if (x is end marker)
        break;
      do something with x;
    }
  }
}
```

# Use In GUIs

- Jave runtime library uses a special thread that does nothing but wait for *events* like mouse clicks, pressed keys, mouse movement, etc.

- You can designate an object of your choice as a *listener;* which means that Java's event thread calls a method of that object whenever an event occurs.

- As a result, your program can do work while the GUI continues to respond to buttons, menus, etc.

- Another special thread does all the drawing. You don't have to be aware when this takes place; just ask that the thread wake up whenever you change something.

# Highlights of a GUI Component

```java
/** A widget that draws multi-colored lines indicated by mouse. */
class Lines extends JComponent implements MouseListener {
  private List<Point> lines = new ArrayList<Point>();

  Lines() {  // Main thread calls this to create one
    setPreferredSize(new Dimension(400, 400));
    addMouseListener(this);
  }
  public synchronized void paintComponent(Graphics g) { // Paint thread
    g.setColor(Color.white);   g.fillRect(0, 0, 400, 400);
    int x, y;    x = y = 200;
    Color c = Color.black;
    for (Point p : lines)
      g.setColor(c); c = chooseNextColor(c);
      g.drawLine(x, y, p.x, p.y); x = p.x; y = p.y;
    }
  }
  public synchronized void mouseClicked(MouseEvent e) // Event thread
    { lines.add(new Point(e.getX(), e.getY())); repaint(); }
  ...
}
```

# Interrupts

- An *interrupt* is an event that disrupts the normal flow of control of a program.

- In many systems, interrupts can be totally *asynchronous*, occurring at arbitrary points in a program, the Java developers considered this unwise; arranged that interrupts would occur only at controlled points.

- In Java programs, one thread can interrupt another to inform it that something unusual needs attention:

  ```
  otherThread.interrupt();
  ```

- But `otherThread` does not receive the interrupt until it waits: methods `wait`, `sleep` (wait for a period of time), `join` (wait for thread to terminate), and mailbox deposit and receive.

- Interrupt causes these methods to throw `InterruptedException`, so typical use is like this:

  ```java
  try {
      msg = inBox.receive();
  } catch (InterruptedException e) { HandleEmergency(); }
  ```

# Remote Mailboxes (A Side Excursion)

- RMI: Remote Method Interface allows one program to refer to objects in another program.

- We use it to allow mailboxes in one program be received from or deposited into in another.

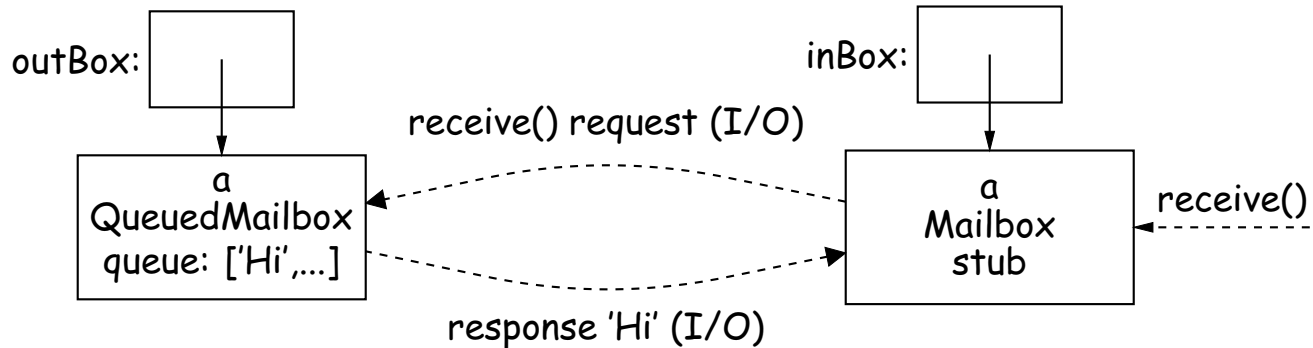- To use this, you define an *interface* to the remote object:

```java
import java.rmi.*;
interface Mailbox extends Remote {
  void deposit(Object msg)
    throws InterruptedException, RemoteException;
  Object receive()
    throws InterruptedException, RemoteException;
  ...
}
```

- On machine that actually will contain the object, you define

```java
class QueuedMailbox ... implements Mailbox {
    Same implementation as before, roughly
}
```

# Remote Objects Under the Hood

```
// On machine #1:              // On Machine #2:
Mailbox outBox                  Mailbox inBox
   = new QueuedMailbox();         = get outBox from machine #1
```

```
outBox: [ ]
         a
      QueuedMailbox  ←----- receive() request (I/O) -----  a
      queue: ['Hi',...]  ----- response 'Hi' (I/O) ----→  Mailbox  ←---- receive()
                                                          stub
inBox: [ ]
```

- Because `Mailbox` is an interface, hides fact that on Machine #2 doesn't actually have direct access to it.

- Requests for method calls are relayed by I/O to machine that has real object.

- Any argument or return type OK if it also implements `Remote` or can be *serialized*—turned into stream of bytes and back, as can primitive types and String.

- Because I/O involved, expect failures, hence every method can throw RemoteException (subtype of IOException).